

Instruction set of the ULM (Ulm Lecture Machine)



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Chapter 1

Description of the ULM

1.1 Data Types

Binary digits are called *bits* and have the value 0 or 1. A *bit pattern* is a sequence of bits. For example

$$X := x_{n-1} \dots x_0 \text{ with } x_k \in \{0, 1\} \text{ for } 0 \leq k < n$$

denotes a bit pattern X with n bits. The number of bits in bit pattern is also called its size or width. The ULM architecture defines a *byte* as a bit pattern with 8 bits. Table 1.1 lists ULM's definitions for *word*, *long word*, *quad word* that refer to specific sizes of bit patterns.

1.2 Expressing the Interpretation of a Bit Pattern

For a bit pattern $X = x_{n-1} \dots x_0$ its *unsigned integer* value is expressed and defined through

$$u(X) = u(x_{n-1} \dots x_0) := \sum_{k=0}^{n-1} x_k \cdot 2^k$$

Signed integer values are represented using the *two's complement* and in this respect the notation

$$s(X) = s(x_{n-1} x_{n-2} \dots x_0) := \begin{cases} u(x_{n-2} \dots x_0), & \text{if } x_{n-1} = 0, \\ u(x_{n-2} \dots x_0) - 2^{n-1}, & \text{else} \end{cases}$$

is used.

1.3 Registers and Virtual Memory

The ULM has 256 registers denoted as %0x00, ..., %0xFF. Each of these registers has a width of 64 bits. The %0x00 is a special purpose register and also denoted as *zero register*. Reading from the zero register always gives a bit pattern where all bits have value 0 (zero bit pattern). Writing to the zero register has no effect.

The (virtual) memory of the ULM is an array of 2^{64} memory cells. Each memory cell can store exactly one byte. Each memory cell has an index which is called its *address*. The address is in the range from 0 to $2^{64}-1$ and the first memory cell of the array has address 0. In notations $M_1(a)$ denotes the memory cell with address a .

Data Size	Size in Bytes	Size in Number of Bits
Bytes	-	8
Word	2	16
Long Word	4	32
Quad Word	8	64

Table 1.1: Names for specific sizes of bit patterns.

1.3.1 Endianness

For referring to data in memory in quantities of words, long words and quad words the definitions

$$\begin{aligned} M_2(a) &:= M_1(a)M_1(a+1) \\ M_4(a) &:= M_2(a)M_2(a+2) \\ M_8(a) &:= M_4(a)M_4(a+4) \end{aligned}$$

are used. The ULM architecture is a *big endian* machine. Therefore we have the equalities

$$\begin{aligned} u(M_2(a)) &= u(M_1(a)M_1(a+1)) \\ u(M_4(a)) &= u(M_2(a)M_2(a+2)) \\ u(M_8(a)) &= u(M_4(a)M_4(a+4)) \end{aligned}$$

1.3.2 Alignment of Data

A quantity of k bytes are aligned in memory if they are stored at an address which is a multiple of k , i.e.

$$M_k(a) \text{ is aligned} \Leftrightarrow a \bmod k = 0$$

Chapter 2

Directives

2.1 .align <expr>

Pad the location counter (in the current segment) to a multiple of <expr>.

2.2 .bss

Set current segment to the BSS segment.

2.3 .byte <expr>

Expression is assembled into next byte.

2.4 .data

Set current segment to the data segment.

2.5 .equ <ident>, <expr>

Updates the symbol table. Sets the value of <ident> to <expr>.

2.6 .global <ident>

Updates the symbol table. Makes the symbol <ident> visible to the linker.

2.7 .globl <ident>

Equivalent to *.globl <ident>*:

Updates the symbol table. Makes the symbol <ident> visible to the linker.

2.8 .long <expr>

Expression <expr> is assembled into next long word (4 bytes).

2.9 .space <expr>

Emits <expr> bytes. Each byte with value 0x00.

2.10 .string <string-literal>

Emits bytes for the zero-terminated <string-literal>.

2.11 .text

Set current segment to the text segment.

2.12 .word <expr>

Expression <expr> is assembled into next word (2 bytes).

2.13 .quad <expr>

Expression <expr> is assembled into next quad word (8 bytes).

Chapter 3

Instructions

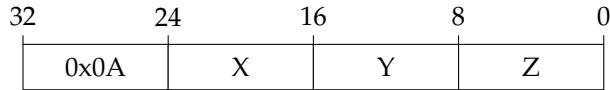
3.1 addq

Integer addition.

3.1.1 Assembly Notation

addq X, %Y, %Z

Format



Effect

$$(u(\%Y) + u(X)) \bmod 2^{64} \rightarrow u(\%Z)$$

Updates the status flags:

Flag	Condition
ZF	$u(\%Y) + u(X) = 0$
CF	$u(\%Y) + u(X) \geq 2^{64}$
OF	$s(\%Y) + s(X) \notin \{-2^{63}, \dots, 2^{63}\}$
SF	$s(\%Y) + s(X) < 0$

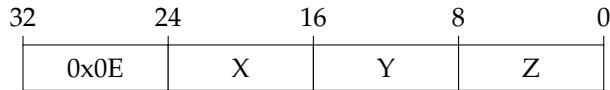
3.1.2 Assembly Notation

addq %X, %Y, %Z

Alternative Assembly Notation

movq %X, %Z

Format



Effect

$$(u(\%Y) + u(\%X)) \bmod 2^{64} \rightarrow u(\%Z)$$

Updates the status flags:

Flag	Condition
ZF	$u(\%Y) + u(\%X) = 0$
CF	$u(\%Y) + u(\%X) \geq 2^{64}$
OF	$s(\%Y) + s(\%X) \notin \{-2^{63}, \dots, 2^{63}\}$
SF	$s(\%Y) + s(\%X) < 0$

3.2 divq

64-bit unsigned integer division

3.2.1 Assembly Notation

divq X, %Y, %Z

Purpose

Uses the 128-bit unsigned integer divider. The 64-bits stored in %Y are therefore zero extended with %0 to 128 bits and then used as numerator. The immediate value X is used as unsigned denominator.

The result of the operation is the quotient and the remainder. These results will be stored in a register pair: The quotient is stored in %Z and the remainder in %{Z + 1}.

Format

32	24	16	8	0
0x10	X	Y	Z	

Effect

For the unsigned 128-bit numerator

$$b := u(\%0\%Y)$$

and unsigned 64-bit denominator

$$a := u(X)$$

computes the divisor

$$\left\lfloor \frac{b}{a} \right\rfloor \bmod 2^{128} \rightarrow u(\%0\%Z)$$

and the remainder

$$b \bmod a \rightarrow u(\%u(Z) + 1)$$

3.2.2 Assembly Notation

divq %X, %Y, %Z

Purpose

Uses the 128-bit unsigned integer divider. The 64-bits stored in %Y are therefore zero extended with %0 to 128 bits and then used as numerator. The register value %X is used as 64-bit unsigned denominator.

The result of the operation is the quotient and the remainder. These results will be stored in a register pair: The quotient is stored in %Z and the remainder in %{Z + 1}.

Format

32	24	16	8	0
0x1A	X	Y	Z	

Effect

For the unsigned 128-bit numerator

$$b := u(\%0\%Y)$$

and unsigned 64-bit denominator

$$a := u(\%X)$$

computes the divisor

$$\left\lfloor \frac{b}{a} \right\rfloor \bmod 2^{128} \rightarrow u(\%0\%Z)$$

and the remainder

$$b \bmod a \rightarrow u(\%\{u(Z) + 1\})$$

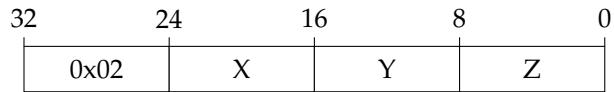
3.3 getc

Get character

3.3.1 Assembly Notation

getc %X

Format



Effect

$$s(ulm_readChar()) \wedge_b 255 \bmod 2^{64} \rightarrow u(\%X)$$

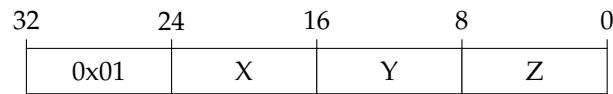
3.4 halt

Halt program

3.4.1 Assembly Notation

halt %X

Format



Effect

halt program execution with exit code $u(\%X) \bmod 2^8$

3.5 imulq

64-bit unsigned and signed integer multiplication.

3.5.1 Assembly Notation

imulq %X, %Y, %Z

Format

32	24	16	8	0
0x0B	X	Y	Z	

Effect

$$u(\%X) \cdot u(\%Y) \bmod 2^{64} \rightarrow u(\%Z)$$

3.5.2 Assembly Notation

imulq X, %Y, %Z

Format

32	24	16	8	0
0x0F	X	Y	Z	

Effect

$$u(X) \cdot u(\%Y) \bmod 2^{64} \rightarrow u(\%Z)$$

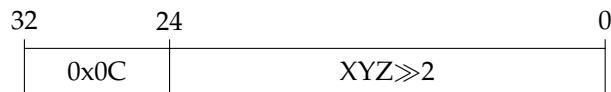
3.6 ja

Jump if above (conditional jump).

3.6.1 Assembly Notation

ja XYZ

Format



Effect

If the condition

$$ZF = 0 \wedge CF = 0$$

evaluates to true then

$$(u(\%IP) + s(XYZ)) \bmod 2^{64} \rightarrow u(\%IP)$$

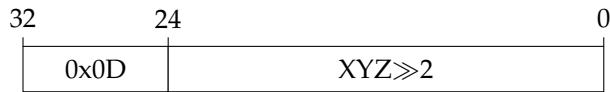
3.7 jb

Jump if below (conditional jump).

3.7.1 Assembly Notation

jb XYZ

Format



Effect

If the condition

$$CF = 1$$

evaluates to true then

$$(u(\%IP) + s(XYZ)) \bmod 2^{64} \rightarrow u(\%IP)$$

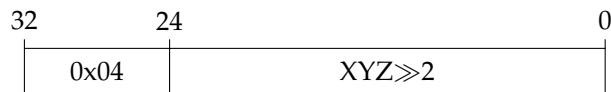
3.8 jmp

Jump (unconditional jump).

3.8.1 Assembly Notation

jmp XYZ

Format



Effect

$$(u(\%IP) + s(XYZ)) \bmod 2^{64} \rightarrow u(\%IP)$$

3.8.2 Assembly Notation

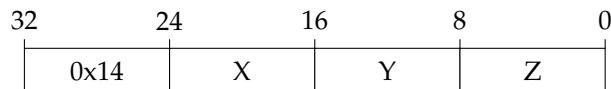
jmp %X, %Y

Alternative Assembly Notation

call %X, %Y

ret %X

Format



Effect

$$\begin{aligned} (u(\%IP) + 4) \bmod 2^{64} &\rightarrow u(\%Y) \\ u(\%X) &\rightarrow u(\%IP) \end{aligned}$$

3.9 jnz

Jump if not zero (conditional jump).

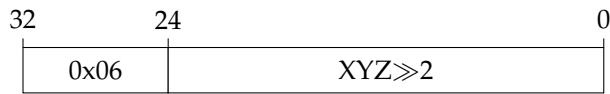
3.9.1 Assembly Notation

jnz XYZ

Alternative Assembly Notation

jne XYZ

Format



Effect

If the condition

$$ZF = 0$$

evaluates to true then

$$(u(\%IP) + s(XYZ)) \bmod 2^{64} \rightarrow u(\%IP)$$

3.10 jz

Jump if zero (conditional jump).

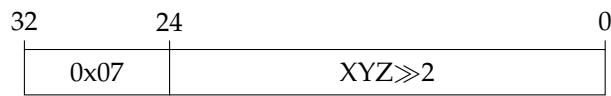
3.10.1 Assembly Notation

jz XYZ

Alternative Assembly Notation

je XYZ

Format



Effect

If the condition

$$ZF = 1$$

evaluates to true then

$$(u(\%IP) + s(XYZ)) \bmod 2^{64} \rightarrow u(\%IP)$$

3.11 ldfp

Load from Pool

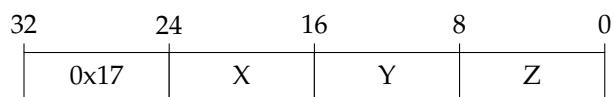
3.11.1 Assembly Notation

ldfp Y(%X), %Z

Alternative Assembly Notation

ldfp (%X), %Z

Format



Effect

$u(M_8(addr)) \rightarrow u(%Z)$ with $addr = (u(Y) \cdot 8 + u(%X)) \bmod 2^{64}$

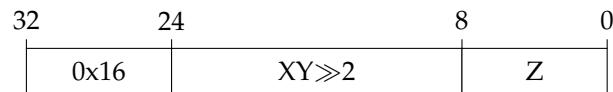
3.12 ldpa

Load Pool Address

3.12.1 Assembly Notation

ldpa XY, %Z

Format



Effect

$$(u(\%IP) + s(XY)) \bmod 2^{64} \rightarrow u(\%Z)$$

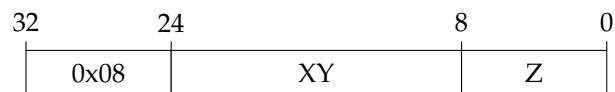
3.13 ldzwq

Load zero extended word to quad word register.

3.13.1 Assembly Notation

ldzwq XY, %Z

Format



Effect

$$u(XY) \bmod 2^{64} \rightarrow u(\%Z)$$

3.14 movb

Move a byte from a register to memory (store instruction).

3.14.1 Assembly Notation

`movb %X, Y(%Z)`

Alternative Assembly Notation

`movb %X, (%Z)`

Purpose

The least significant byte in `%X` gets stored at address `%Z`.

Format

32	24	16	8	0
0x11	X	Y	Z	

Effect

$$u(\%X) \bmod 2^8 \rightarrow u(M_1(addr)) \text{ with } addr = (s(Y) + u(\%Z)) \bmod 2^{64}$$

3.15 movq

Move quad word

3.15.1 Assembly Notation

`movq Y(%X), %Z`

Alternative Assembly Notation

`movq (%X), %Z`

Purpose

Fetches a quad word from address `%X` into register `%Z`

Format

32	24	16	8	0
0x12	X	Y	Z	

Effect

$u(M_8(addr)) \rightarrow u(%Z)$ with $addr = (s(Y) + u(%X)) \bmod 2^{64}$

3.15.2 Assembly Notation

`movq %X, Y(%Z)`

Alternative Assembly Notation

`movq %X, (%Z)`

Format

32	24	16	8	0
0x19	X	Y	Z	

Effect

$u(%X) \bmod 2^{64} \rightarrow u(M_8(addr))$ with $addr = (s(Y) + u(%Z)) \bmod 2^{64}$

3.16 movzbq

Move zero extended byte to quad word register (fetch instruction).

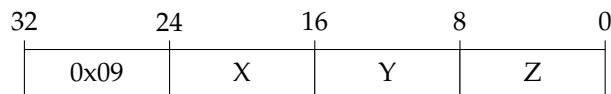
3.16.1 Assembly Notation

movzbq Y(%X), %Z

Alternative Assembly Notation

movzbq (%X), %Z

Format



Effect

$$u(M_1(addr)) \rightarrow u(%Z) \text{ with } addr = (s(Y) + u(%X)) \bmod 2^{64}$$

3.17 putc

Put character

3.17.1 Assembly Notation

putc %X

Format

32	24	16	8	0
0x03	X	Y	Z	

Effect

$$\text{ulm_printChar}(u(\%X) \wedge_b 255)$$

3.17.2 Assembly Notation

putc X

Format

32	24	16	8	0
0x13	X	Y	Z	

Effect

$$\text{ulm_printChar}(u(X) \wedge_b 255)$$

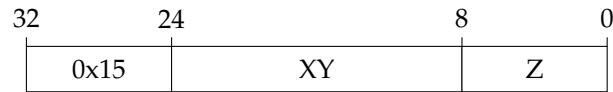
3.18 shldwq

Shift (destination register) and load word

3.18.1 Assembly Notation

shldwq XY, %Z

Format



Effect

$$u(\%Z) << 16 | u(XY) \bmod 2^{64} \rightarrow u(\%Z)$$

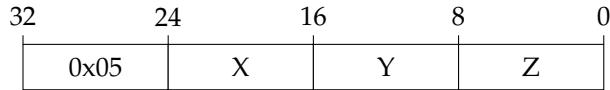
3.19 subq

Integer subtraction.

3.19.1 Assembly Notation

subq X, %Y, %Z

Format



Effect

$$(u(\%Y) - u(X)) \bmod 2^{64} \rightarrow u(\%Z)$$

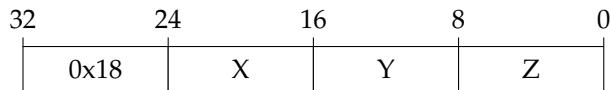
Updates the status flags:

Flag	Condition
ZF	$u(\%Y) - u(X) = 0$
CF	$u(\%Y) - u(X) < 0$
OF	$s(\%Y) - s(X) \notin \{-2^{63}, \dots, 2^{63}\}$
SF	$s(\%Y) - s(X) < 0$

3.19.2 Assembly Notation

subq %X, %Y, %Z

Format



Effect

$$(u(\%Y) - u(\%X)) \bmod 2^{64} \rightarrow u(\%Z)$$

Updates the status flags:

Flag	Condition
ZF	$u(\%Y) - u(\%X) = 0$
CF	$u(\%Y) - u(\%X) < 0$
OF	$s(\%Y) - s(\%X) \notin \{-2^{63}, \dots, 2^{63}\}$
SF	$s(\%Y) - s(\%X) < 0$

Chapter 4

ISA Source File for the ULM Generator

```
1  RRR (OP u 8) (X u 8) (Y u 8) (Z u 8)
2  RSR (OP u 8) (X u 8) (Y s 8) (Z u 8)
3  J26 (OP u 8) (XYZ j 24)
4  UI16R (OP u 8) (XY u 16) (Z u 8)
5  J18R (OP u 8) (XY j 16) (Z u 8)
6
7
8  0x01 RRR
9  : halt %X
10   ulm_halt(ulm_regVal(X));
11
12 0x02 RRR
13 : getc %X
14   ulm_setReg(ulm_readChar() & 0xFF, X);
15
16 0x03 RRR
17 : putc %X
18   ulm_printChar(ulm_regVal(X) & 0xFF);
19
20 0x04 J26
21 : jmp XYZ
22   ulm_unconditionalRelJump(XYZ);
23
24 0x05 RRR
25 : subq X, %Y, %Z
26   ulm_sub64(X, ulm_regVal(Y), Z);
27
28 0x06 J26
29 : jnz XYZ
30 : jne XYZ
31   ulm_conditionalRelJump(ulm_statusReg[ULM_ZF] == 0, XYZ);
32
33 0x07 J26
34 : jz XYZ
35 : je XYZ
```

```

36     ulm_conditionalRelJump(ulm_statusReg[ULM_ZF] == 1, XYZ);
37
38 0x08 U16R
39 : ldzwq XY, %Z
40     ulm_setReg(XY, Z);
41
42 0x09 RSR
43 : movzbq Y(%X), %Z
44 : movzbq (%X), %Z
45     ulm_fetch64(Y, X, 0, 0, ULM_ZERO_EXT, 1, Z);
46
47 0x0A RRR
48 : addq X, %Y, %Z
49     ulm_add64(X, ulm_regVal(Y), Z);
50
51 0x0B RRR
52 : imulq %X, %Y, %Z
53     ulm_mul64(ulm_regVal(X), ulm_regVal(Y), Z);
54
55 0x0C J26
56 : ja XYZ
57     ulm_conditionalRelJump(ulm_statusReg[ULM_ZF] == 0
58                               && ulm_statusReg[ULM_CF] == 0, XYZ);
59
60 0x0D J26
61 : jb XYZ
62     ulm_conditionalRelJump(ulm_statusReg[ULM_CF] == 1, XYZ);
63
64 0x0E RRR
65 : addq %X, %Y, %Z
66 : movq %X, %Z
67     ulm_add64(ulm_regVal(X), ulm_regVal(Y), Z);
68
69 0x0F RRR
70 : imulq X, %Y, %Z
71     ulm_mul64(X, ulm_regVal(Y), Z);
72
73 0x10 RRR
74 # Uses the 128-bit unsigned integer divider. The 64-bits stored in %Y are
75 # therefore zero extended with %0 to 128 bits and then used as numerator.
76 # The immediate value X is used as unsigned denominator.
77 #
78 # The result of the operation is the quotient and the remainder. These results
79 # will be stored in a register pair: The quotient is stored in %Z and the
80 # remainder in %{Z + 1}.
81 : divq X, %Y, %Z
82     ulm_div128(X, ulm_regVal(Y), ulm_regVal(0), Z, 0, Z + 1);
83
84 0x11 RSR
85 # The least significant byte in %X gets stored at address %Z.
86 : movb %X, Y(%Z)
87 : movb %X, (%Z)
88     ulm_store64(Y, Z, 0, 0, 1, X);

```

```

89
90 0x12 RSR
91 # Fetches a quad word from address %X into register %Z
92 : movq Y(%X), %Z
93 : movq (%X), %Z
94     ulm_fetch64(Y, X, 0, 0, ULM_ZERO_EXT, 8, Z);
95
96 0x13 RRR
97 : putc X
98     ulm_printChar(X & 0xFF);
99
100 # void ulm_absJump(uint64_t destAddr, ulm_Reg ret);
101 0x14 RRR
102 : jmp %X, %Y
103 : call %X, %Y
104 : ret %X
105     ulm_absJump(ulm_regVal(X), Y);
106
107 0x15 U16R
108 : shldwq XY, %Z
109     ulm_setReg(ulm_regVal(Z) << 16 | XY, Z);
110
111 0x16 J18R
112 : ldpa XY, %Z
113     ulm_setReg(ulm_ipVal() + XY, Z);
114
115 0x17 RRR
116 : ldfp Y(%X), %Z
117 : ldfp (%X), %Z
118     ulm_fetch64(Y * 8, X, 0, 0, ULM_ZERO_EXT, 8, Z);
119
120 0x18 RRR
121 : subq %X, %Y, %Z
122     ulm_sub64(ulm_regVal(X), ulm_regVal(Y), Z);
123
124 0x19 RSR
125 : movq %X, Y(%Z)
126 : movq %X, (%Z)
127     ulm_store64(Y, Z, 0, 0, 8, X);
128
129 0x1A RRR
130 # Uses the 128-bit unsigned integer divider. The 64-bits stored in %Y are
131 # therefore zero extended with %0 to 128 bits and then used as numerator.
132 # The register value %X is used as 64-bit unsigned denominator.
133 #
134 # The result of the operation is the quotient and the remainder. These results
135 # will be stored in a register pair: The quotient is stored in %Z and the
136 # remainder in %{Z + 1}.
137 : divq %X, %Y, %Z
138     ulm_div128(ulm_regVal(X), ulm_regVal(Y), ulm_regVal(0), Z, 0, Z + 1);
139
140 # General meaning of the mnemonics

```

```
142  
143 @addq  
144 # Integer addition.  
145 @getc  
146 # Get character  
147 @divq  
148 # 64-bit unsigned integer division  
149 @halt  
150 # Halt program  
151 @imulq  
152 # 64-bit unsigned and signed integer multiplication.  
153 @ja  
154 # Jump if above (conditional jump).  
155 @jb  
156 # Jump if below (conditional jump).  
157 @je  
158 # Jump if equal (conditional jump).  
159 @jmp  
160 # Jump (unconditional jump).  
161 @jne  
162 # Jump if not equal (conditional jump).  
163 @jnz  
164 # Jump if not zero (conditional jump).  
165 @jz  
166 # Jump if zero (conditional jump).  
167 @ldzwq  
168 # Load zero extended word to quad word register.  
169 @putc  
170 # Put character  
171 @subq  
172 # Integer subtraction.  
173 @movb  
174 # Move a byte from a register to memory (store instruction).  
175 @movzbq  
176 # Move zero extended byte to quad word register (fetch instruction).  
177 @movq  
178 # Move quad word  
179 @shldwq  
180 # Shift (destination register) and load word  
181 @ldpa  
182 # Load Pool Address  
183 @ldfp  
184 # Load from Pool
```
